

Plaid Parliament of Pwning 2024 eCTF Team Carnegie Mellon University

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Presentation Outline

- Design Phase
 - Overview
 - Design Highlight: Zero-Trust Architecture

Attack Phase

- I2CBleed Exploit
- Supply Chain I2CBleed
- Other Attacks + Interesting Defenses
- Project Management + Lessons Learned



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Our Design Highlights

Custom Hardened Encrypted Link Encryption At-Rest of Everything **Physical Link Layer** Layer Wrapper **ChaCha-Poly AEAD Random Nonces to** Board RNG + vonfor encryption **Prevent Replays** Neumann **Minimal External Avoid Interrupts &** Random Delays + **Code Surface** Async Code **Redundant Checks**

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Design Highlight: Zero-Trust Architecture

- **Thought Experiment:** Assume full hardware compromise
 - How to defend flags? Can we use fun crypto tricks?
- BB Boot / BB Extract: Encrypt comp. secrets w/ key stored in AP
- Op. PIN Extract / SC Extract: Encrypt keys inside AP w/ PIN
 - *Potential for offline brute-force if AP compromised
- Op. Pump Swap: Not defensible, but encrypt the code to make it harder
- SC Boot / Damaged Boot: ?????
 - How to require both components to be present in order to boot?







Design Highlight: Zero-Trust Architecture

- Damaged Boot: Require all components be present in order to boot?
- "Russian Encryption Doll": Encrypt AP boot data with all component keys
- How to distribute component keys?
 - Comp Key = Hash(Root Key || Comp ID)
- How to do replace component?
 - Keep Root Key encrypted with Replace Token
 - RT is long enough to not be brute-able

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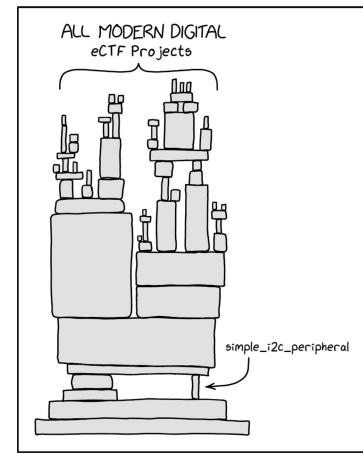
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Attack Highlight: I2CBleed

- Three vulnerabilities in starter code
 - Read/write indices not reset on repeated start
 - Read index checked for == instead of >=
 - Write index casted to unsigned (overflows)
- Result: Arbitrary Read/Write (!!!) (of anything past I2C_REGS)
- Straightforward Attack Process
 - 1. Write in malicious shellcode
 - 2. Write a bunch of padding
 - 3. Overwrite vector table to jump to shellcode



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Attack Highlight: Fully-Automated I2CBleed



- Q: What to do with a near-universal arbitrary-code-execution exploit?
 - A: Make it full auto: <u>4-5 flags in 90 seconds from ZIP download</u>
- Step 1: Determine I2C Address of Victim
 - Scan all addresses, see which ones ACK (like insecure list_components)
- Step 2: Determine I2C_REGS address (shellcode address)
 - Use arbitrary read until the component crashes (stops ACKing)
- Step 3: Inject shellcode
 - Step 3.5 (SC only): Scan until we find the string "ctf{"
 - Locally: Dump all of flash to the UART (including keys and plaintext flags!!)
- Step 4 (SC only): Bitbang SPI data back to malicious component
 - Malicious component receives SPI and dumps anything transmitted over UART





Interesting Defenses

- Defending against I2CBleed
 - Certificate Chain: Provide each component with a ID-unique certificate signed using a deployment-time CA
 - Encrypt component attestation data / boot message with key stored in AP
 - Key pinning to assign unique component keys (bypass deployment hash check...)
- Other unique defenses
 - Challenge-response handshake on every message in the system
 - Custom I2C implementation (don't trust provided libraries...)
 - Use of hardware features / PUFs to prevent emulation

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Project Management + Lessons Learned

• Design Phase

- Get everyone set up with insecure example in the first week
- Design security protocol *before* starting implementation, but can start generic tasks (scripting, infra, comms, crypto library) simultaneously
- Secure By Design: Drive out the attacker in every possible way

Attack Phase

- Balance between optimizing conventional attacks and developing novel attacks
- Track red-team availability for executing rapid attacks for first bloods
- Be willing to operate at strange hours (sadly)

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Project Management + Lessons Learned

- Overall
 - Earning course credit helps offset the time investment
 - Cross-Training: EEs studied crypto, Security students studied electronics
 - If viable, hardware setup for each team member to individually play with

Lessons Learned

- Sustainability of having most of the work be done by a few team members?
- Redundancy to avoid single points of failure (esp. for design phase timeline)
- Novel attacks require *a lot* more human-hours than estimated, fine-tuning "standard" attacks can be better

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Thank you!



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Extra Slides





